# Integrated Circuits for Digital Communications 

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## Basic Baseband PAM Concepts

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## General Data Communication System



- Source coder removes redundancy from source (i.e. MPEG, ADPCM, text compression, etc.)
- Channel coder introduces redundancy to maximize information rate over channel.
(i.e. error-correcting codes, trellis coding, etc.)
- Our interest is in channel coding/decoding and channel transmission/reception.


## Basic Baseband System



- In 2B1Q, coder maps pairs of bits to one of four levels - $A_{k}=\{-3,-1,1,3\}$


## Rectangular Transmit Filter



- The spectrum of $A_{k}$ is flat if random.
- The spectrum of $s(t)$ is same shape as $H_{t}(f)$


## Nyquist Pulses

- $h(t)$ is the impulse response for transmit filter, channel and receive filter ( $\otimes$ denotes convolution)

$$
\begin{gather*}
h(t)=h_{t}(t) \otimes h_{c}(t) \otimes h_{r}(t)  \tag{1}\\
q(t)=\sum_{m=-\infty}^{\infty} A_{m} h(t-m T)+n(t) \otimes h_{r}(t)
\end{gather*}
$$

- The received signal, $q(t)$, is sampled at $k T$.

$$
q_{k}=\sum_{m=-\infty}^{\infty} A_{m} h(k T-m T)+u(k T) \quad, u(t) \equiv n(t) \otimes h_{r}(t)
$$

- For zero intersymbol interference (i.e. $q_{k}=A_{k}+u_{k}$ )

$$
\begin{equation*}
h(k T)=\delta_{k} \quad\left(\delta_{k}=0,1,0,0,0, \ldots\right) \tag{4}
\end{equation*}
$$

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## Nyquist Pulses

- For zero ISI, the same criteria in the frequency domain is: $\left(f_{s}=1 / T\right)$

$$
\frac{1}{T} \sum^{\infty} H\left(j 2 \pi f+j m 2 \pi f_{s}\right)=1
$$

- Known as Nyquist Criterion

Example Nyquist Pulses (in freq domain)


Sinc pulse



## Nyquist Pulses



Sinc pulse


Raised-cosine pulse


Sinc pulse


Raised-cosine pulse

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## Raised-Cosine Pulse



- $\alpha$ determines excess bandwidth


## Raised-Cosine Pulses



- More excess bandwidth - impulse decays faster.

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## Raised-Cosine Pulse

- $\alpha$ determines amount of excess bandwidth past $f_{s} / 2$
- Example: $\alpha=0.25$ implies that bandwidth is 25 percent higher than $f_{s} / 2$ while $\alpha=1$ implies bandwidth extends up to $f_{s}$.
- Larger excess bandwidth - easier receiver
- Less excess bandwidth - more efficient channel use


## Example

- Max symbol-rate if a $50 \%$ excess bandwidth is used and bandwidth is limited to 10 kHz
- $1.5 \times\left(f_{s} / 2\right)=10 \mathrm{kHz}$ implies $f_{s}=13.333 \times 10^{3}$ symbols $/ \mathrm{s}$


## Eye Diagram



- "a" indicates immunity to noise
- "b" indicates immunity to errors in timing phase
- slope "c" indicates sensitivity to jitter in timing phase


## Eye Diagram

- Zero crossing - NOT a good performance indicator
- $100 \%$ bandwidth has little zero crossing jitter
- $50 \%$ BW has alot of zero crossing jitter but it is using less bandwidth


- Less excess BW - more intolerant to timing phase


## Example Eye Diagrams


$\alpha=0.550 \%$ excess bandwidth


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## Example Eye Diagrams


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## Matched-Filter



- For zero-ISI, $h_{t c}(t) \otimes h_{r}(t)$ satisfies Nyquist criterion.
- For optimum noise performance, $h_{r}(t)$ should be a matched-filter.
- A matched-filter has an impulse response which is time-reversed of $h_{t c}(t)$

$$
\begin{equation*}
h_{r}(t)=K h_{t c}(-t) \tag{6}
\end{equation*}
$$

where $K$ is an arbitrary constant.

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## Matched-Filter (proof)

- Consider isolated pulse case (so no worry about ISI)

$$
\begin{gather*}
r(t)=A_{0} h_{t c}(t)+n(t)  \tag{7}\\
q_{0}=\left.\int_{-\infty}^{\infty} r(\tau) h_{r}(t-\tau) d \tau\right|_{t=0}=\int_{-\infty}^{\infty} r(\tau) h_{r}(-\tau) d \tau  \tag{8}\\
q_{0}=A_{0}^{\infty} \int_{-\infty}^{\infty} h_{t c}(\tau) h_{r}(-\tau) d \tau+\int_{-\infty}^{\infty} n(\tau) h_{r}(-\tau) d \tau \tag{9}
\end{gather*}
$$

- Want to maximize signal term to noise term
- Variance of noise is

$$
\begin{equation*}
\sigma_{n}^{2}=N_{0} \int_{-\infty}^{\infty} h_{r}^{2}(-\tau) d \tau \tag{10}
\end{equation*}
$$

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## Matched-Filter (proof)

- Assuming $A_{0}$ and $h_{t c}(t)$ fixed, want to maximize

$$
\begin{equation*}
\mathrm{SNR}=\frac{A_{0}^{2}\left[\int_{-\infty}^{\infty} h_{t c}(\tau) h_{r}(-\tau) d \tau\right]^{2}}{N_{0}^{2} \int_{-\infty}^{\infty} h_{r}^{2}(-\tau) d \tau} \tag{11}
\end{equation*}
$$

- Use Schwarz inequality

$$
\begin{equation*}
\left[\int_{a}^{b} f_{1}(x) f_{2}(x) d x\right]^{2} \leq\left[\int_{a}^{b} f_{1}^{2}(x) d x\right]\left[\int_{a}^{b} f_{2}^{2}(x) d x\right] \tag{12}
\end{equation*}
$$

with equality if and only if $f_{2}(x)=K f_{1}(x)$

- Maximizing (11) results in $h_{r}(t)=K h_{t c}(-t)$ - QED


## Matched-Filter - Why optimum?



Transmit filter, channel and noise



Too much noise,
All of signal


## ISI and Noise

- In general, we need the output of a matched filter to obey Nyquist criterion
- Frequency response at output of matched filter is $\left|H_{t c}(j \omega)\right|^{2}$ leading to criterion

$$
\begin{equation*}
\frac{1}{T} \sum_{m=-\infty}^{\infty}\left|H_{t c}\left(j 2 \pi f+j m 2 \pi f_{s}\right)\right|^{2}=1 \tag{13}
\end{equation*}
$$

## Example

- Assume a flat freq resp channel and raised-cosine pulse is desired at matched-filter output
- Transmit filter should be $\sqrt{\text { raised-cosine }}$
- Receive filter should be $\sqrt{\text { raised-cosine }}$


# Gaussian Noise and SNR Requirement 

## Probability Distribution Function

- Consider a random variable X
- Cumulative distribution function (c.d.f.) $-F_{x}(x)$

$$
\begin{gather*}
F_{x}(x) \equiv P_{r}(X \leq x)-\infty<x<\infty  \tag{14}\\
1 \geq F_{x}(x) \geq 0 \tag{15}
\end{gather*}
$$

Example

- Consider a fair die


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## Probability Density Function

- Derivative of $F_{x}(x)$ is p.d.f. defined as $f_{x}(x)$

$$
\begin{equation*}
f_{x}(x) \equiv \frac{d F_{x}(x)}{d x} \quad \text { or } \quad F_{x}(x)=\int_{-\infty}^{\alpha} f_{x}(\alpha) d \alpha \tag{16}
\end{equation*}
$$

- To find prob that $X$ is between $x_{1}$ and $x_{2}$

$$
\begin{equation*}
P_{r}\left(x_{1}<X \leq x_{2}\right)=\int_{x}^{x_{2}} f_{x}(\alpha) d \alpha \tag{17}
\end{equation*}
$$

- It is the area under p.d.f. curve.

Example (fair die)


## Uniform Distribution

- p.d.f. is a constant
- Variance is given by: $\sigma^{2}=\frac{\Delta^{2}}{12}$ where $\Delta$ is range of random variables

- Crest factor: $C F \equiv \frac{\max }{\sigma}=\frac{\Delta / 2}{\Delta / \sqrt{12}}=\sqrt{3}=1.732$


## Example

- A uniform random variable chosen between 0 and 1 has a mean, $\mu=0.5$, and variance, $\sigma^{2}=1 / 12$

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## Gaussian Random Variables

## Probability Density Function

- Assuming $\sigma^{2}=1$ (i.e. variance is unity) and $\mu=0$ (i.e. mean is zero) then

$$
\begin{equation*}
f_{x}(x)=\frac{1}{\sqrt{2 \pi}} e^{-x^{2} / 2} \tag{18}
\end{equation*}
$$



## Gaussian Random Variables

- Often interested in how likely a random variable will be in tail of a Gaussian distribution

$$
\begin{gather*}
Q(x) \equiv P_{r}(X>x)=\frac{1}{\sqrt{2 \pi}} \int_{x}^{\infty} e^{-\alpha^{2} / 2} d \alpha  \tag{19}\\
Q(x)=\frac{1}{2} \operatorname{erfc}(x / \sqrt{2}) \tag{20}
\end{gather*}
$$



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## Gaussian Random Variables

- Probability of $x$ being in tail of Gaussian distribution

- If $\sigma^{2} \neq 1$ or $\mu \neq 0$

$$
\begin{equation*}
P_{r}(X>x)=Q((x-\mu) / \sigma) \tag{21}
\end{equation*}
$$

## Example SNR Calculation

- 100Base-T2 for fast-ethernet uses 5-PAM
- Want to calculate the receive SNR needed for a symbol-error-rate of $10^{-10}$ (assume rest is ideal).

- Signal power, $P_{s}$

$$
\begin{equation*}
P_{s}=\frac{1}{4} \times 0 W+\frac{1}{2} \times 4 W+\frac{1}{4} \times 16 W=6 W \tag{22}
\end{equation*}
$$

- Using a reference of $1 W$ as 0 dB ,

$$
\begin{equation*}
P_{s}=10 \log _{10}(6)=7.78 \mathrm{~dB} \tag{23}
\end{equation*}
$$

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## Example SNR Calculation

- Assume Gaussian noise added to receive signal.
- Since symbols are distance 2 apart, a noise value greater than 1 will cause an error in receive symbol.
- Want to find $\sigma$ of Gaussian distribution such that likelihood of random variable greater than 1 is $10^{-10}$.
- Recall

$$
\begin{equation*}
Q(x / \sigma)=0.5 \operatorname{erfc}((x / \sigma) / \sqrt{2}) \tag{24}
\end{equation*}
$$

- Let $x=1$ and set

$$
\begin{equation*}
2 Q(1 / \sigma)=10^{-10} \tag{25}
\end{equation*}
$$

( 2 value because variable might be $>1$ or $<-1$ )

$$
\begin{equation*}
0.5 \times 10^{-10}=Q(1 / \sigma)=0.5 \operatorname{erfc}(1 /(\sigma \sqrt{2})) \tag{26}
\end{equation*}
$$

## Example SNR Calculation

- Trial and error gives $1 /(\sigma \sqrt{2})=4.57$ implying that $\sigma=0.1547=1 / 6.46$
- Noise with $\sigma=0.1547$ has a power of (ref to $1 W$ )

$$
\begin{equation*}
P_{n}=10 \log _{10}\left(\sigma^{2}\right)=-16.2 \mathrm{~dB} \tag{27}
\end{equation*}
$$

- Finally, SNR needed at receive signal is

$$
\begin{equation*}
\mathrm{SNR}=7.78 \mathrm{~dB}-(-16.2 \mathrm{~dB})=24 \mathrm{~dB} \tag{28}
\end{equation*}
$$

- Does not account that large positive noise on +4 signal will not cause symbol error (same on -4).
- It is slightly conservative
- BER approx same as symbol error rate if Gray coded


## m-PAM

- For $m$ bits/symbol $\Rightarrow 2^{m}$ levels
- Normalize distance between levels to 2 (so error of 1 causes a symbol error)
- $(m=1) \Rightarrow \pm 1 \quad(m=3) \Rightarrow \pm 1, \pm 3, \pm 5, \pm 7 \quad$ etc.
- Noise variance of $(\sigma=0.1547) \Rightarrow B E R=10^{-10}$
- Symbols spaced $\pm 1, \pm 3, \pm 5, \ldots, \pm\left(2^{m}-1\right)$
- average power is: $S_{m}=\left(4^{m}-1\right) / 3$

$$
\begin{equation*}
\mathrm{SNR}=10 \log \left(\frac{S_{m}}{\sigma^{2}}\right)=10 \log \left(\frac{4^{m}-1}{3 \sigma^{2}}\right) \tag{29}
\end{equation*}
$$

## m-PAM

$$
\begin{equation*}
\mathrm{SNR}=10 \log \left(\frac{S_{m}}{\sigma^{2}}\right)=10 \log \left(\frac{4^{m}-1}{3 \sigma^{2}}\right) \tag{30}
\end{equation*}
$$

- equals 23.1 dB for $m=2$, $\mathrm{BER}=10^{-10}$
- equals 28.2 dB for $m=3$, $\mathrm{BER}=10^{-10}$ (approx +6 dB )
- Can show $S_{m+1}=4 S_{m}+1$
- Require 4 times more power to maintain same symbol error rate with same noise power (uncoded)
- In other words,
- to send 1 more bit/symbol, need 6dB more SNR (but does not increase bandwidth)


## Why Assume Gaussian Noise?

## Central-Limit Theorem

- Justification for modelling many random signals as having a Gaussian distribution

> Sum of independent random variables approaches Gaussian as sum increases

- Assumes random variables have identical distributions.
- No restrictions on original distribution (except finite mean and variance).
- Sum of Gaussian random variables is also Gaussian.


## Uniform and Gaussian Signals

1000 samples of uniform random variables


1000 samples of Gaussian random variables


## Filtered Random Signals



Filtered with 3'rd order Butterworth lowpass with cutoff $f_{s} / 200$

No longer independent from sample to sample


## Wired Digital Communications

## Wired Digital Transmission

Long Twisted-Pair Applications ( $1 \mathbf{k m}$ - 6 km )

- T1/E1 - $1.5 / 2 \mathrm{Mb} / \mathrm{s}(2 \mathrm{~km})$
- ISDN - Integrated Services Digital Network
- HDSL - High data-rate Digital Subscriber Line
- ADSL - Asymmetric DSL
- VDSL - Very high data-rate DSL

Short Twisted-Pair Applications (20m - 100m)

- $100 \mathrm{Mb} / \mathrm{s}$ Fast-Ethernet - TX, T4, T2
- Gigabit Ethernet - Short haul, Long haul


## Short Coax (300m)

- Digital video delivery - 300Mb/s - $1.5 \mathrm{~Gb} / \mathrm{s}$


## Cable Modelling

- Modelled as a transmission line.



## Twisted-Pair Typical Parameters:

- $R(f)=(1+j) \sqrt{f / 4} \Omega / \mathrm{km} \quad$ due to the skin effect
- $L=0.6 \mathrm{mH} / \mathrm{km}$ (relatively constant above 100 kHz )
- $C=0.05 \mu \mathrm{~F} / \mathrm{km}$ (relatively constant above 100 kHz )
- $G=0$


## Skin Effect

- "Resistance" is not constant with frequency and is complex valued.
- Can be modelled as:

$$
\begin{equation*}
R(\omega)=k_{R}(1+j) \sqrt{\omega} \tag{31}
\end{equation*}
$$

where $k_{R}$ is a constant given by

$$
\begin{equation*}
k_{R}=\frac{1}{\pi d_{c}}\left(\frac{\mu}{2 \sigma}\right)^{1 / 2} \tag{32}
\end{equation*}
$$

- $d_{c}$ is conductor diameter, $\mu$ is permeability, $\sigma$ is conductivity
- Note resistance is inversely proportional to $d_{c}$.
- Jordan and Balmain, "Electromagnetic Waves and Radiating Systems", pg. 563, Prentice-Hall, 1968.


## Characteristic Impedance

$$
\begin{equation*}
Z_{0}=\sqrt{\frac{R+j \omega L}{G+j \omega C}} \tag{33}
\end{equation*}
$$

- Making use of (31) and assuming $G=0$

$$
\begin{align*}
& Z_{0}=\left(\frac{k_{R} \sqrt{\omega}(1+j)+j \omega L}{j \omega C}\right)^{1 / 2}  \tag{34}\\
& Z_{0}=\sqrt{\frac{L}{C}}\left(1+\frac{k_{R}}{L \sqrt{\omega}}(1-j)\right)^{1 / 2} \tag{3}
\end{align*}
$$

Now using approx $(1+x)^{1 / 2} \approx 1+x / 2$ for $x \ll 1$

$$
\begin{equation*}
Z_{0} \approx \sqrt{\frac{L}{C}}+\frac{k_{R}}{2 \sqrt{\omega L C}}(1-j) \tag{3}
\end{equation*}
$$

- At high freq, $Z_{0}$ appears as constant value $\sqrt{L / C}$


## Characteristic Impedance

- From (33), when $\omega L » R$ (typically $\omega » 2 \pi \times 16 \mathrm{kHz}$ )

$$
\begin{equation*}
Z_{0 \mathrm{~h}}=\sqrt{\frac{L}{C}} \tag{37}
\end{equation*}
$$

resulting in

$$
\begin{equation*}
Z_{\mathrm{Oh}} \approx 110 \Omega \tag{38}
\end{equation*}
$$

- Thus, when terminating a line, a resistance value around $110 \Omega$ should be used.


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## Cable Transfer-Function

- When properly terminated, a cable of length $d$ has a transfer-function of

$$
\begin{equation*}
H(d, \omega)=e^{-d \gamma(\omega)} \tag{3}
\end{equation*}
$$

where $\gamma(\omega)$ is given by

$$
\begin{equation*}
\gamma(\omega)=\sqrt{(R+j \omega L)(G+j \omega C)} \tag{40}
\end{equation*}
$$

- Breaking $\gamma(\omega)$ into real and imaginary parts,

$$
\begin{gather*}
\gamma(\omega) \equiv \alpha(\omega)+j \beta(\omega)  \tag{41}\\
H(d, \omega)=e^{-d \alpha(\omega)} e^{-j d \beta(\omega)} \tag{42}
\end{gather*}
$$

- $\alpha(\omega)$ determines attenuation.
- $\beta(\omega)$ determines phase.


## Cable Transfer-Function

- Assuming $G=0$, then from (40)

$$
\begin{equation*}
\gamma=\left(j \omega C R-\omega^{2} L C\right)^{1 / 2} \tag{43}
\end{equation*}
$$

- Substituting in (31)

$$
\begin{align*}
& \gamma=\left(j \omega^{1.5} k_{R} C(1+j)-\omega^{2} L C\right)^{1 / 2}  \tag{44}\\
& \gamma=j \omega \sqrt{L C}\left(1+\frac{k_{R}}{L \sqrt{\omega}}(1-j)\right)^{1 / 2} \tag{45}
\end{align*}
$$

Now using approx $(1+x)^{1 / 2} \approx 1+x / 2$ for $x \ll 1$

$$
\begin{equation*}
\gamma \approx \frac{k_{R}}{2} \sqrt{\frac{\omega C}{L}}+j\left(\omega \sqrt{L C}+\frac{k_{R}}{2} \sqrt{\frac{\omega C}{L}}\right) \tag{46}
\end{equation*}
$$

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## Cable Attenuation

- Equating (41) and (46)

$$
\begin{equation*}
\alpha(\omega) \approx \frac{k_{R}}{2} \sqrt{\frac{C}{L}} \times \sqrt{\omega} \tag{47}
\end{equation*}
$$

- Therefore gain in dB is

$$
\begin{equation*}
H_{d B}(d, \omega) \approx-8.68 d \times \frac{k_{R}}{2} \sqrt{\frac{C}{L}} \times \sqrt{\omega} \tag{48}
\end{equation*}
$$

- Note that attenuation in dB is proportional to cable length (i.e. $2 x$ distance doubles attenuation in dB)
- Can reduce attenuation by using a larger diameter cable
- Attenuation proportional to root-frequency


## Cable Attenuation

- Gain in dB is proportional to $\sqrt{f}$ due to skin effect.

- Do not confuse with $1 / f$ noise slow frequency roll-off.

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## Cable Phase

- Equating (41) and (46)

$$
\begin{equation*}
\beta(\omega) \approx \omega \sqrt{L C}+\frac{k_{R}}{2} \sqrt{\frac{C}{L}} \times \sqrt{\omega} \tag{49}
\end{equation*}
$$

- The linear term usually dominates
- The linear term implies a constant group delay.
- In other words, the linear term simply accounts for the delay through the cable.
- Ignoring linear phase portion, remaining phase is proportional to $\sqrt{ } f$.
- Note it has the same multiplying term as attenuation.


## IIR Filter Cable Match using Matlab

\% this program calculates an iir num/den transfer-function
\% approx for a transmission line with exp(sqrt(s)) type response.
clear;
\% Order of IIR filter to match to cable
$\% \mathrm{nz}$ is numerator order and np is denominator order
nz = 9;
$n p=10 ;$
\% important parameters of cable
$\mathrm{c}=0.05 \mathrm{e}-6 \%$ capacitance per unit length in farads $/ \mathrm{km}$
$\mathrm{I}=0.6 \mathrm{e}-3 \%$ inductance per unit length in henries $/ \mathrm{km}$
$\mathrm{kr}=0.25 \%$ resistance per unit length in ohms/km (times ( $1+\mathrm{j})^{*}$ sqrt(omega))
$\mathrm{d}=0.1 \%$ cable length in km
\% above values adjusted to obtain -20dB atten for 100 m at 125 MHz
k_cable $=(\mathrm{kr} / 2)^{*}$ sqrt(c/l);
\% the frequency range for finding tf of cable fmin=1;
fmax=1e9;
\% specify frequency points to deal with nmax=1000;
$\mathrm{f}=\operatorname{logspace}(\log 10(\mathrm{fmin}), \log 10(\mathrm{fmax}), \mathrm{nmax})$;
$\mathrm{w}=2^{*} \mathrm{p} \mathrm{i}^{\star} \mathrm{f}$;
$s=j^{*} w$;
\% 'cable' is desired outcome in exponential form cable $=\exp \left(-d^{*} k \_\right.$cable ${ }^{*}$ sqrt(2) ${ }^{*}$ sqrt(s));
\% Perform IIR approximate transfer-function match
\% Since invfreqs miminizes (num-cable*den)
\% first need an approximate den so that it can be used
$\%$ as a freq weighting to minimize (num/den - cable)
[num,den]=invfreqs(cable,w,nz,np, 1./w);
[denor]=freqs(den,1,w);
\% re-iterate process with weighting for the denominator
\% which now minimizes (num/den - cable)
[num,den]=invfreqs(cable,w,nz,np, (1./denor). ${ }^{\wedge} 2$ );
[denor]=freqs(den,1,w);
[num,den]=invfreqs(cable,w,nz,np, (1./denor). ${ }^{\wedge} 2$ );
\% find approximate transfer function 'cable_approx' to 'cable' [cable approx]=freqs(num,den,w);
\% also find pole-zero model
[Z,p,k]=tf2zp(num,den);

## \% PLOT RESULTS

clf;
figure(1);
subplot(211);
semilogx(f,20*log10(abs(cable)), ${ }^{\prime}{ }^{\prime}$ ');
hold on;
semilogx(f,20*log10(abs(cable_approx)),'b');
title('Cable Magnitude Response’);
xlabel('Freq (Hz)');
ylabel('Gain (dB)’);
grid;
hold off;
subplot(212);
semilogx(f,angle(cable)*180/pi,'r');
hold on;
semilogx(f,angle(cable_approx)*180/pi,'b');
title('Cable Phase Response');
xlabel ('Freq (Hz)');


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ylabel('Phase (degrees)');
grid;
hold off;
figure(2);
subplot(211);
semilogx(f,20*log10(abs(cable)./abs(cable_approx)));
title('Gain Error Between Cable and Cable_approx');
xlabel('Freq (Hz)');
ylabel('Gain Error (dB)');
subplot(212);
semilogx(f,(angle(cable)-angle(cable_approx))*180/pi); title('Phase Error Between Cable and Cable_approx');
xlabel('Freq (Hz)');
ylabel('Phase Error (degrees)’);
grid;

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## Cable Response



Cable Phase Response


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## IIR Matching Results



Phase Error Between Cable and Cable ${ }_{a}$ pprox


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## Near and Far End Crosstalk



- In FEXT, interferer and signal both attenuated by cable
- In NEXT, signal attenuated but interferer is coupled directly in.
- When present, NEXT almost always dominates.
- Can cancel NEXT if nearby interferer is known.
- Envelope of squared gain of NEXT increases with $f^{1.5}$


## Twisted-Pair Crosstalk

- Crosstalk depends on turns/unit length, insulator, etc.
- Twisted-pairs should have different turns/unit length within same bundle


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## Transformer Coupling

- Almost all long wired channels (>10m) are AC coupled systems
- AC coupling introduces baseline wander if random PAM sent
- A long string of like symbols (for example, +1) will decay towards zero degrading performance
- Requires baseline wander correction (non-trival)
- Can use passband modulation schemes (CAP, QAM, DMT)
- Why AC couple long wired channels??


## Transformer Coupling

## Eliminates need for similar grounds

- If ground potentials not same - large ground currents
Rejects common-mode signals
- Transformer output only responds to differential signal current
- Insensitive to common-mode signal on both wires



## Generic Wired PAM Transceiver



- Look at approaches for each block


## HDSL Application

- $1.544 \mathrm{Mb} / \mathrm{s}$ over 4.0 km of existing telephone cables.
- Presently 4-level PAM code (2B1Q) over 2 pairs (a CAP implementation also exists).
- Symbol-rate is $386 \mathrm{ksymbols} / \mathrm{s}$


## Possible Bridged-Taps

- Can have unterminated taps on line
- Modelling becomes more complicated but DFE equalizes effectively
- Also causes a wide variation in input line impedance to which echo canceller must adapt - difficult to get much analog echo cancellation

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## HDSL Application

- Symbol-rate is 386 ksymbols/s


## Received Signal

- For $d=4 k m$, a 200 kHz signal is attenuated by $40 d B$.
- Thus, high-freq portion of a 5 Vpp signal is received as a 50 mV pp signal - Need effective echo cancellation


## Transmit Path

- Due to large load variations, echo cancellation of analog hybrid is only 6dB
- To maintain 40dB SNR receive signal, linearity and noise of transmit path should be better than 74 dB .


## ISDN Application

- Similar difficulty to HDSL but lower frequency
- $160 \mathrm{~kb} / \mathrm{s}$ over 6 km of 1 pair existing telephone cables
- 4-level PAM coding - 2B1Q
- Receive signal at 40 kHz atten by 40 dB
- Requires highly linear line-drivers + A/D converters for echo cancellation (similar to HDSL)


## Fast-Ethernet Application

| CAT3 | CAT5 |
| :---: | :---: |
| $H_{d B}(f)=2.32 \sqrt{f}+0.238 f$ | $H_{d B}(f)=1.967 \sqrt{f}+0.023 f+0.05 / \sqrt{f}$ |
| $12.5 M H z \leftrightarrow 11 d B$ | $12.5 M H z \leftrightarrow 7 d B$ |
| crosstalk worse | crosstalk better |

## 100Base-T4

- 4 pair CAT3 - 3 pair each way, 25MS/s with coding 100Base-TX
- 2 pair CAT5 - 3 level PAM to reduce radiation 100Base-T2
- 2 pair CAT3 - $5 \times 5$ code, $25 \mathrm{MS} / \mathrm{s}$ on each pair


## Typical Transmit D/A Block



- Polyphase filter to perform upsampling+filtering


## HDSL

- D/A and filter needs better than 12-bit linearity
- Might be an oversampled 1-bit DAC
- One example: $\boldsymbol{\wedge}_{16} ; 48$ tap FIR; $\boldsymbol{\uparrow}_{4} ; \Delta \Sigma$ DAC


## Fast-Ethernet

- Typically around 35 dB linearity + noise requirement
- 100Base-T2 example: $\uparrow$ 3 ; simple FIR; 75 MHz 4 -bit DAC; 3'rd-order LP cont-time filter


## Line Drivers

- Line driver supplies drive current to cable.
- Commonly realized as voltage buffers.
- Often the most challenging part of analog design.
- Turns ratio of transformer determines equivalent line impedance.

$V_{n e}=\frac{2}{n} V_{2}$

$$
\begin{gathered}
V_{1}=V_{2} / n \\
I_{1}=n I_{2} \\
R_{1}=R_{2} / n^{2}
\end{gathered}
$$

Typical Values
$R_{2}=100 \Omega$
$V_{2}= \pm 2.5 \mathrm{~V}$
$I_{2}= \pm 25 \mathrm{~mA}$

## Line Driver Efficiency

- Efficiency improves as power supply increased


## Example (assume can drive within 1 V of supplies)

- From typical values, max power delivered by line driver is $P_{\text {linetR }}=2 \times 2.5 \times 25 \mathrm{~mA}=125 \mathrm{~mW}$


## 12V Case

- Consider 12V supply - use $n=0.5, V_{n e, \max }=10 \mathrm{~V}$, $I_{1, \text { max }}=12.5 \mathrm{~mA}$ leading to $P=12 \times 12.5 \mathrm{~mA}=150 \mathrm{~mW}$ (and drive an 800 ohm load)


## 3V Case

- Consider 3V supply - use $n=5, V_{n e, \text { max }}=1 \mathrm{~V}$, $I_{1, \text { max }}=125 \mathrm{~mA}$ leading to $P=3 \times 125 \mathrm{~mA}=375 \mathrm{~mW}$ (and drive an 8 ohm load!!!)


## Line Driver

- In CMOS, W/L of output stage might have transistors on the order of 10,000!
- Large sizes needed to ensure some gain in final stage so that feedback can improve linearity - might be driving a 30 ohm load
- When designing, ensure that enough phase margin is used for the wide variation of bias currents
- Nested Miller compensation has been successfully used in HDSL application with class AB output stage
- Design difficulties will increase as power supplies decreased


## 2-4 Wire Hybrids

- Dual-duplex often used to reduce emission.
- However, dual-duplex requires hybrids and echo cancellation.

- If $R_{L}=R_{T}$, no echo through hybrid
- Can be large impedance variation.

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## Typical HDSL Line Impedances



## Hybrid Issues

- Note zero at dc and pole at 10 kHz .
- Low frequency pole causes long echo tail (HDSL requires 120 tap FIR filter)


## Alternatives

- Could eliminate $R_{1}$ circuit and rely on digital echo cancellation but more bits in A/D required.
OR
- Can make $R_{1}$ circuit more complex to ease A/D specs.
- Less echo return eases transmit linearity spec.
- Might be a trend towards active hybrids with or without extra A/D and D/A converters (particularly for higher speeds).


## Typical Receive A/D



- Often, VGA is controlled from digital signal.
- Anti-aliasing can be simple in oversampled systems.
- Continuous-time filters are likely for fast-ethernet
- Example: 100Base-T2 suggests a 5'th order conttime filter at 20 MHz with a $6-$ bit $\mathrm{A} / \mathrm{D}$ at 75 MHz .
- Challenge here is to keep size and power of $A / D$ small.


## Echo Cancellation



- Typically realized as an adaptive FIR filter.
- Note input is transmit signal so delay lines and multiplies are trivial.
- HDSL uses about a 120 tap FIR filter
- Coefficient accuracy might be around 20 bits for dynamic range of 13 bits.


## Echo Cancellation

- Fast-ethernet might be around 30 taps and smaller coefficient accuracy
- Can also perform some NEXT cancellation if signal of nearby transmitter is available (likely in 100Base-T2 and gigabit ethernet)


## Alternatives

- Higher data rates may have longer echo tails.
- Might go to FIR/IIR hybrid to reduce complexity.
- Non-linear echo cancellation would be VERY useful in reducing transmit linearity spec.
- However, these non-linearities have memory and thus Volterra series expansions needed.


## Equalization

## HDSL

- Echo canceller required before equalization so fractional spaced equalizer not practical
- Typically 9 tap FFE and 120 tap DFE
- Long DFE also performs dc recovery (baseline wander)
Fast Ethernet
- Often fractional-spaced EQ - 30 taps
- DFE - 20 taps (dc recovery)

