

CSC 2228 Project

Dynamic Voltage Scaling in Mobile Devices

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What We Did?

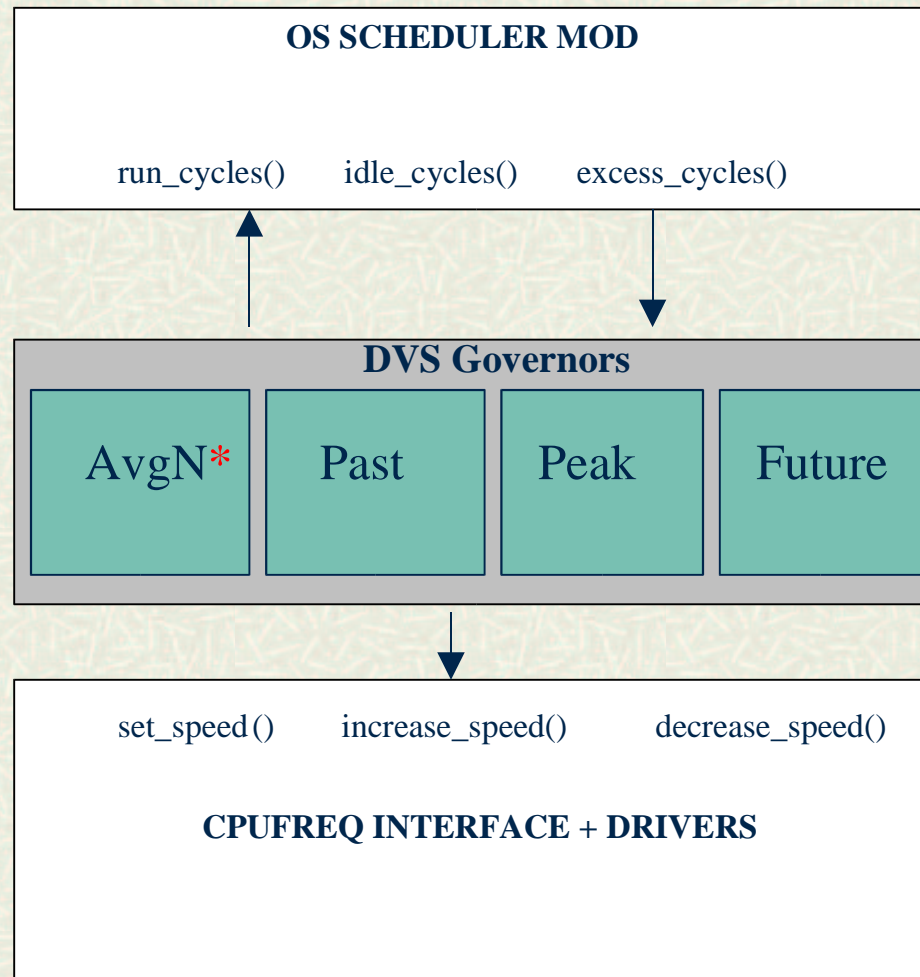
- # Implement DVS algs
+ develop our own
 - # Target: Linux 2.6 & Intel P4-M
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Why is this important?

- # DVS goal = save energy
 - # Explore DVS algs on Linux 2.6 & P4-M SpeedStep
 - # Used hypothetical idle times
 - # Developed new DVS alg with OS info
 - # DVS alg:
 - AvgN*, PAST, PEAK, FUTURE
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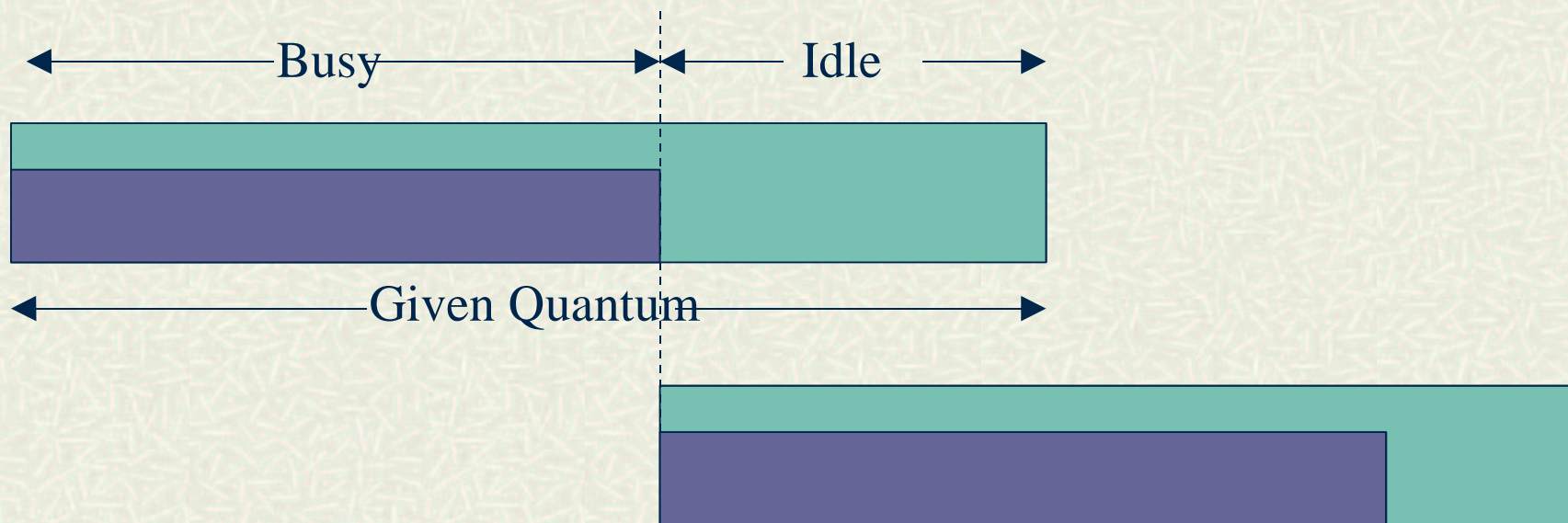
Architecture

3 Components:



DVS Interpretation

- # Goal: stretch task to consume 99% of the given quantum



- # No major changes to scheduler

DVS ALGORITHMS

FUTURE

- Our own proposed algorithm
- Idea:
 - Look into the run_queue:
 - Get hypothetical idle time in the future
- Prediction and Speed Setting is same as PAST

EXPERIMENTAL SETUP

Platform:

- IBM Thinkpad T30
- Redhat 9 2.6.0-test9 Linux kernel.
- Intel Mobile Pentium 4 - M processor

Frequency (GHz)	Voltage (V)
1.6	1.3
1.2	1.2
Idle-state	-

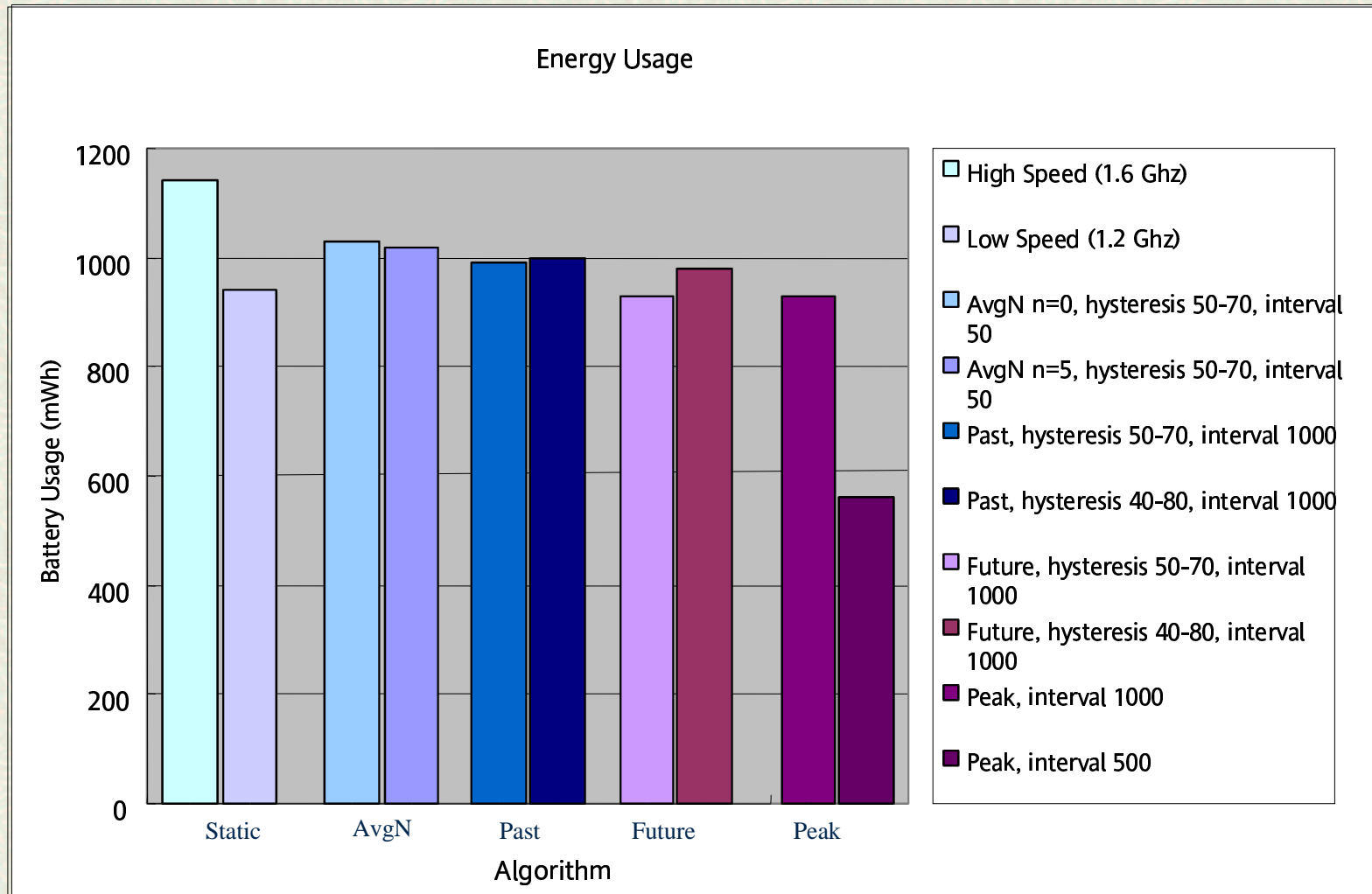
- Used ACPI in Linux to measure capacity (mWh)

EXPERIMENTAL SETUP

Experiments:

- Workload is 2m30sec long and consists of:
 - An mpeg video running in the background
 - Sleep for 20 seconds
 - Compile Linux kernel (linking stage)
 - lasts 20-30 seconds
 - Sleep for 20 seconds
 - Same compilation again
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EXPERIMENTAL RESULTS



FUTURE WORK

- # Implementation of excess cycle
 - # Fine tuning DVS alg parameters
 - # More workloads
 - # New DVS algs
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CONCLUSION & CONTRIBUTIONS

- # Best policies are AVGn and PAST
- # DVS algs on new OS & real hardware
 - Used hypervisor idle times
- # “Future” alg
- # Contributions to Linux community:
 - 4 governors for CPUFreq
 - Make available on web